

• Linear Search :-

The simplest of all the searching is linear or sequential search. Sequential searching is nothing but searching element in linear way, we have a need to start the search from beginning and search the element one by one until the end of array or linked list. If search is successful then it will return the location otherwise it will return the failure notification.

eg. →

	0	1	2	3	4
	30	60	90	12	80

 Key element → 12

Comparing Key element with first element of array.

1 $a[0] == \text{Key}$
 $30 == 12$ X

2 $a[1] == \text{Key}$
 $60 == 12$ X

3 $a[2] == \text{Key}$
 $90 == 12$ X

4

$a[3] = \text{Key}$
 $12 = 12$

search successful element
 found at position 4.

Q

WAP to search element in an
 array using linear search
 technique.

```
#include <stdio.h>
#include <conio.h>
void main ()
{
    int a[10], n, Key, flag = 0, i;
    clrscr ();
    printf ("Enter how many element  

    you want in array");
    scanf ("%d", &n);
    printf ("In Enter element");
    for (i = 0; i < n; i++)
    {
        scanf ("%d", &a[i]);
    }
    printf ("Element of array");
    {
        printf ("%d", a[i]);
    }
    printf ("Enter element to be  

    searched");
```

```
Search (" %d", &key)
```

```
{  
  if (a[i] == key)
```

```
  {  
    flag = 1;  
    break;
```

```
  }  
}
```

```
if (flag == 1)  
  printf ("Element found at  
  position %d", i+1);
```

```
else
```

```
{  
  printf ("\n Element not found  
  in array");
```

```
}  
getch();  
}
```

Algorithm for Linear Search.

Here we represent an unsorted array of element and n represent number of element and key element represent the value to be searched for in the array.

Step - 1 [Initialize]
 $i = 0$, $flag = 0$

Step - 2 \rightarrow Repeat step III for $i = 0, 1, 2, 3, 4, 5 \dots n$.

Step - 3 if $(a[i] == \text{key})$, then
flag = 1
printf ("Element is found at %d", i+1);
STOP;
end if

Step - 4 if (flag == 0), then
printf ("Element was not found in the array")
end if
stop

Binary Searching

A unsorted array is searched by linear search that scan the array element one by one until the desired element is found.

The reason for sorting the array is that we can search the array. Now if the array is sorted we can employ binary search which brilliantly divide the list into two part. The size of the search space is halved

each time it examine one array element.

An array based binary search selects the middle element in the array and compare its value to that of the key element, because the array is sorted if the key is less than the middle element in the array then the key must be in the first half of the array. Likewise if the value of key element is greater than that of the middle value in the array then the key lies in the second half of the array.

In ideal case we in fact check one half of the search space or array with only one comparison.

So, the algorithm narrows the search area by half at each time until it has either found a key or search fails.

As the name suggests, binary search divide the array into halves. This search is applicable only to ordered list.

Q WAP to search a key element in an array using a binary search technique.

```
#include <stdio.h>
#include <conio.h>
void main()
{
    int a[20], n, i, Key, low, high, mid,
        flag = 1;
    clrscr();
    printf("\n Enter the element you
        want to store in array");
    scanf("%d", &n);
    low = 0;
    high = n - 1;
    printf("Enter the element in
        sorted order.");
    for(i = 0; i < n; i++)
    {
        scanf("%d", &a[i]);
    }
    printf("\n Enter searched Element");
    scanf("%d", &Key);
    while (low <= high)
    {
        mid = (low + high) / 2;
        if (a[mid] == Key)
        {
            flag = 0;
        }
    }
}
```



```
flag = 1;  
break;  
}  
else if (key > a[mid])  
{  
    low = mid + 1;  
}  
else  
{  
    high = mid - 1;  
}  
}  
if (flag == 1)  
{  
    printf("\n Element found at position  
        %d", mid + 1);  
}  
else  
{  
    printf("Element is not found in  
        array or list");  
}  
getch();  
}
```

Difference between linear search and binary search

Linear Search

Binary Search

1 It is sequential type of search

1 It is divide and conquer approach

2 The element can be searched if the list of element are in unsorted order.

2 Necessary and mandatory condition is that the element should be in sorted order i.e ascending or descending order

3 Linear search is efficient if the list contain his element

3 Binary search is efficient if list has large number of elements.

Sorting

i) Bubble Sorting :-

In this ~~sorting~~ algorithm, multiple swapping take place in one pass.

Smaller elements move or 'bubble' up to the top of the list, hence the name given to the algorithm.

In this method, adjacent members of list to be sorted are compared. If the item on top is greater than item immediately below it, then they are swapped. This process is carried on till the list is sorted.

Q We have following N numbers

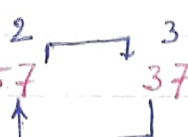
0	1	2	3	4	5	6	7
25	57	48	37	12	92	86	33

Step I ie $A[0]$ with $A[1]$ ie $A[0] < A[1]$ ie $25 < 57$ so no interchange.

$A[1]$ with $A[2]$ ie $A[1] > A[2]$ so interchange

Step II -

0	1	2	3	4	5	6	7
25	48	57	37	12	92	86	33



$A[2] > A[3]$ so

Step III

0	1	2	3	4	5	6	7
25	48	37	57	12	92	86	33

step (IV)

0	1	2	3	4	5	6	7
25	48	37	57	12	92	86	33

$A[3] > A[4]$

so interchange

step (V)

0	1	2	3	4	5	6	7
25	48	37	12	57	92	86	33

$57 < 92$ so not interchange

step (VI)

$92 > 86$ ie $A[5] > A[6]$ so interchange

0	1	2	3	4	5	6	7
25	48	37	12	57	86	92	33

step (VII)

$A[6] > A[7]$ so interchange

25	48	37	12	57	86	33	92
----	----	----	----	----	----	----	----

25	37	48	12	57	86	33	92
----	----	----	----	----	----	----	----

25	37	12	48	57	86	33	92
----	----	----	----	----	----	----	----

25	37	12	48	57	33	86	92
----	----	----	----	----	----	----	----

25	12	37	48	57	33	86	92
----	----	----	----	----	----	----	----

25	12	37	48	33	57	86	92
----	----	----	----	----	----	----	----

12	25	37	48	33	57	86	92
----	----	----	----	----	----	----	----

12	25	37	33	48	57	86	92
----	----	----	----	----	----	----	----

12 | 25 | 33 | 37 | 48 | 57 | 86 | 92 ✓

Algorithm :-

- 1) Begin
- 2) Read the n elements
- 3) for $i = 1$ to n
for $j = n$ down to $i + 1$
if $a[j] < a[j - 1]$
swap ($a[j]$, $a[j - 1]$)
- 4) End

Total No of comparison in Bubble sort

$$= (N-1) + (N-2) + \dots + 2 + 1$$
$$= (N+1) * N / 2 = O(N^2)$$

Efficiency :-

In the best case if all elements are sorted in the array no interchange is made. Under the worst case condition i.e. when there are $n-1$ passes & $n-1$ comparison have to be made on each pass total no of comparison = $(n-1) * (n-2)$ which is $O(n^2)$.

Advantages :-

- 1) Simple to understand
- 2) Easy to implement

Disadvantages :-

- 1) slowest sorting technique
- 2) Most inefficient sorting technique.

Q

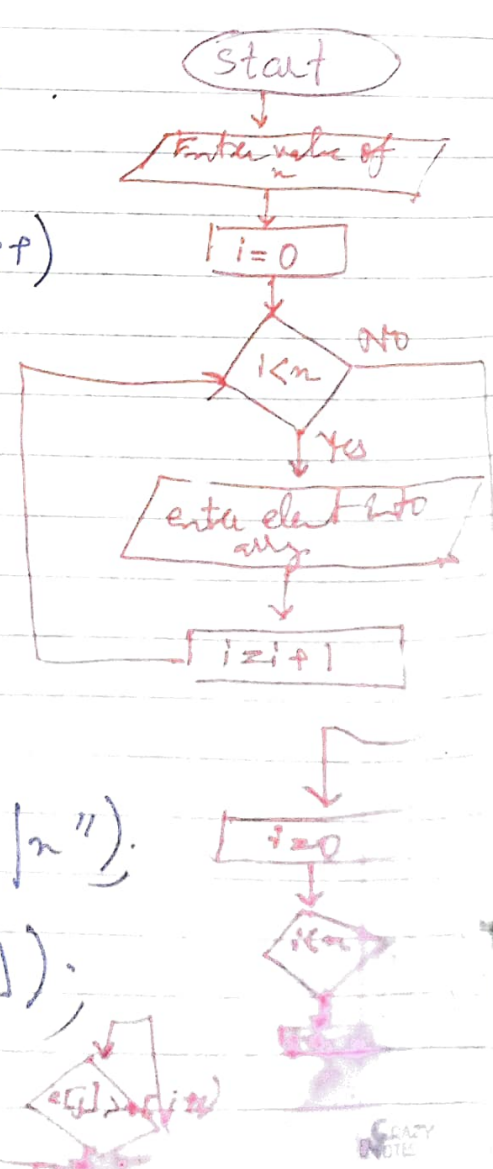
Write a program to sort the elements using Bubble sorting technique.

Sol) =>

```

#include <stdio.h>
#include <conio.h>
void main()
{
    int a[20], n, i, j, temp;
    clrscr();
    printf("\n how many no in array |n");
    scanf("%d", &n);
    printf("\n enter the element into array |n");
    for (i=0; i<n; i++)
    {
        scanf("%d", &a[i]);
    }
    printf("\n Unsorted array |n");
    for (i=0; i<n; i++)
        printf("%d |t", a[i]);
    for (i=0; i<n; i++)
    {
        for (j=0; j<n-1-i; j++)
        {
            if (a[j] > a[j+1])
            {
                temp = a[j];
                a[j] = a[j+1];
                a[j+1] = temp;
            }
        }
    }
    printf("\n Sorted Array |n");
    for (i=0; i<n; i++)
        printf("%d |t", a[i]);
    getch();
}

```



(2) Insertion sort :-

- 1) An insertion sort is one that sort a set of values by inserting values into an existing sorted file.
- 2) The main idea of the algorithm is to build a complete solution by inserting a new element from the unsorted portion of a list.
- 3) In an array a with n element $a[1], a[2], \dots, a[n]$ is in the memory. The insertion sort algorithm scans a from $a[1]$ to $a[n]$ inserting each element $a[k]$ into its proper position in previously sorted subarray.

$a[1], \dots, a[k-1]$

Advantages :-

- (i) Simple to understand
- (ii) Easy to implement

Disadvantages :-

- (i) works inefficiently on large amount of data.

0	1	2	3	4	5	6
25	15	30	9	99	20	26

pass I: - $a[1] < a[0]$ interchanging the position of element.

15	25	30	9	99	20	26
0	1	2	3	4	5	6

pass II: - $a[3]$ is less than $a[0]$, $a[1]$, $a[2]$

9	15	25	30	99	20	26
0	1	2	3	4	5	6

pass III: - $a[4] > a[3]$ so No change

pass (IV): - $a[5]$ is less than $a[2]$, $a[3]$, $a[4]$

9	15	20	25	30	99	26
0	1	2	3	4	5	6

pass (V): - $a[6]$ is less than $a[4]$, $a[5]$, $a[6]$

9	15	20	25	26	30	99
0	1	2	3	4	5	6

Algorithm: -

Step 1: - For $I=1$ to N

Step 2: - $J=I$

Step 3: - While ($j \geq 1$)

Step 4: - If ($A[j] < A[j-1]$) then

Step 5: - $temp = A[j]$

Step 6: - $A[j] = A[j-1]$

Step 7: - $A[j-1] = temp$

Step 8: - End If

$J=J-1$

[End of while Loop]

[End of step 1 For Loop]

pg: Exit

Q Write a program to sort elements of an array using insertion sort technique.

```
sol) => #include <stdio.h>
#include <conio.h>
void main()
{
    int a[100], n;
    int i, j, temp;
    clrscr();
    printf("\n Enter Number of elements ");
    scanf("%d", &n);
    printf("\n Enter element into array (n times);");
    for (i=0; i<n; i++)
        scanf("%d", &a[i]);
    printf("\n The Unsorted list (n times);");

    for (i=0; i<n; i++)
        printf("%d\t", a[i]);
    for (i=1; i<n; i++)
    {
        j=i;
        while (j>=1)
        {
            if (a[j]<a[j-1])
            {
                temp = a[j];
                a[j] = a[j-1];
                a[j-1] = temp;
            }
            j=j-1;
        }
    }
    printf("\n The sorted list is (n times);");
    for (i=0; i<n; i++)
        printf("%d\t", a[i]);
}
```



```
printf("%d\n");
getch();
}
```

Example: -

	0	1	2	3	4	5	6	7
	82	42	49	8	92	25	59	52

Step I: - $a[1] < a[0]$ so interchanging $a[0]$ & $a[1]$

	0	1	2	3	4	5	6	7
	42	82	49	8	92	25	59	52

Step II: - $a[2] < a[1]$ but $a[2] \neq a[0]$ so ^{only} interchanging 1 & 2.

	0	1	2	3	4	5	6	7
	42	49	82	8	92	25	59	52

Step III: - Here $a[3]$ is less than $a[0]$, $a[1]$, $a[2]$

	0	1	2	3	4	5	6	7
	8	42	49	82	92	25	59	52

Step IV: -

	0	1	2	3	4	5	6	7
	8	42	49	82	92	25	59	52

 Here $a[5]$ is less than $a[1]$, $a[2]$, $a[3]$, $a[4]$.

Step V: -

	0	1	2	3	4	5	6	7
	8	25	42	49	82	92	59	52

Step VI: -

	0	1	2	3	4	5	6	7
	8	25	42	49	59	82	92	52

Step VII:

	0	1	2	3	4	5	6	7
	8	25	42	49	52	59	82	92

selection sort

If we have a list of elements in unsorted order & we want to make a list of elements in sorted order then first we will take the smallest element & keep in the new list, after that second smallest element & so on until the biggest element of list.

algo

pass I:

- ① search the smallest element from $arr[0] \dots arr[N-1]$
- ② Interchange $arr[0]$ with smallest element.
Result $arr[0]$ is sorted.

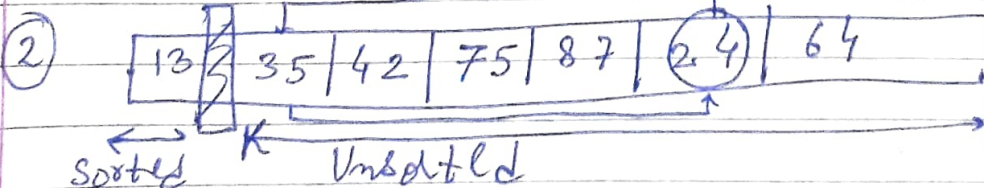
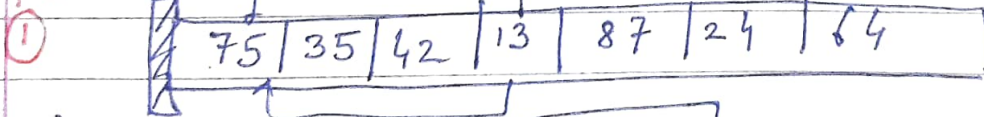
pass II:-

- ① search the smallest element from $arr[1] \dots arr[N-1]$
- ② Interchange $arr[1]$ with smallest element
Result $arr[0], arr[1]$ is sorted

pass N-1:-

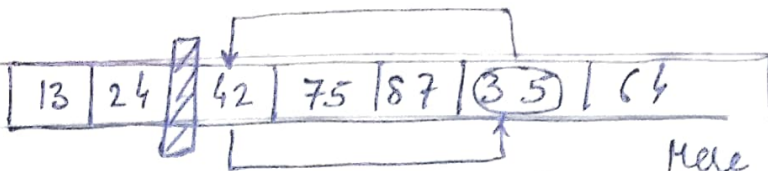
- ① search the smallest element from $arr[N-2] \dots arr[N-1]$.
- ② Interchange $arr[N-2]$ with smallest element.
Result :- $arr[0] \dots arr[N-1]$ is sorted.

pass



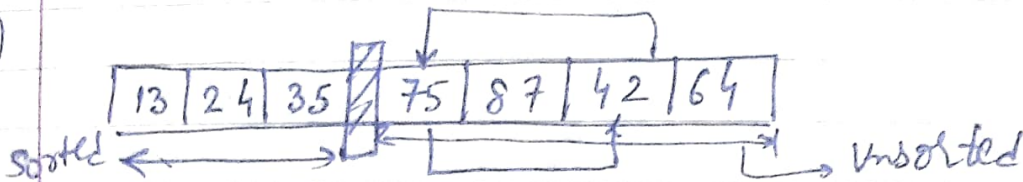
from $arr[1] \dots arr[n-1]$
select

3



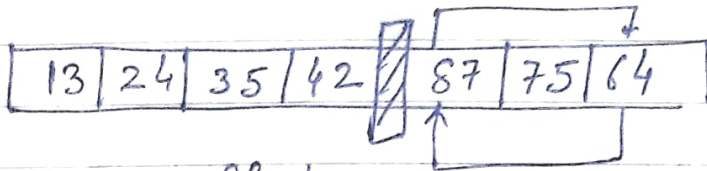
Here smallest is 35.

4



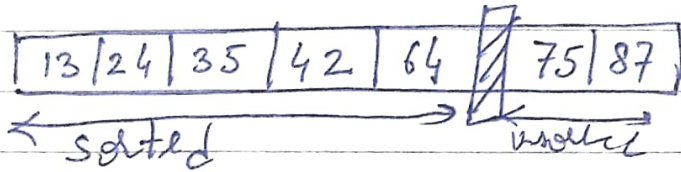
Here smallest element in unsorted part is 42

5



Here smallest in unsorted is 64.

6



output :-

13 24 35 42 64 75 87

Analysis :-

$$F(n) = (n-1) + (n-2) + (n-3) + \dots + 3 + 2 + 1$$

$$\text{Sum} = \frac{n}{2} [2a + (n-1)d]$$

$a \rightarrow$ First element in series

$n =$ NO of element in the series

$d =$ diff between second element & first element
(ie consecutive element) ie 2-1

$$a = (n-1) \quad d \rightarrow \cancel{n} - 2 - \cancel{n} + 1$$

Hence

$$F(n) = \frac{(n-1)}{2} [2(n-1) + \{(n-1) - 1\} \{(n-2) - (n-1)\}]$$

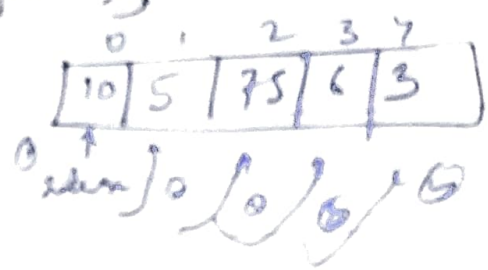
$$= \frac{(n-1)}{2} [2n - 2 + (n-2)(-1)]$$

$$= \frac{(n-1)}{2} [2n - 2 - n + 2]$$

$$= \frac{(n-1)}{2} \times n = \boxed{O(n^2)}$$

Ex Write a program to sort the elements of array using selection sort.

```
#include <stdio.h>
#include <conio.h>
void main()
{
    int arr[20], i, j, k, n, temp, smallest;
    printf("\n Enter the number of elements |n");
    scanf("%d", &n);
    for(i=0; i<n; i++)
    {
        printf("\n Enter element %d:", i+1);
        scanf("%d", &arr[i]);
    }
    printf("\n Unsorted list: |n");
    for(i=0; i<n; i++)
        printf("%d", arr[i]);
    printf("\n");
    for(i=0; i<n-1; i++)
    {
        smallest = i;
        for(k=i+1; k<n; k++)
        {
            if(arr[smallest] > arr[k])
                smallest = k;
        }
        if(i != smallest)
        {
            temp = arr[i];
            arr[i] = arr[smallest];
            arr[smallest] = temp;
        }
    }
}
```



```
printf ("\n After pres %d elements are: ", i+1);  
for (j=0; j<n; j++)  
    printf ("%d", arr[j]);  
    printf ("\n");  
} /* End of for */
```

```
printf ("\n sorted list is: \n");  
for (i=0; i<n; i++)  
    printf ("%d", arr[i]);  
    printf ("\n");  
}
```

Q Write a C++ program to sort the elements of array using quick sort technique

```
#include <iostream.h>
#include <conio.h>
#define max 100
class quicksort
{
    int i, l, h;
public:
    void input();
    void output(int *, int);
    void quick-sort(int *, int, int);
};
void main()
{
    int i, l, h, n, a[max];
    clrscr();
    quicksort qs;
    cout << endl << "How many elements in the array";
    cin >> n;
    cout << endl << "Enter the elements : ";
    for (i = 0; i <= n - 1; i++)
    {
        cin >> a[i];
    }
    l = 0;
    h = n - 1;
    qs.quick-sort(a, l, h);
    cout << "sorted array";
    qs.output(a, n);
    getch();
}
```



```
void quicksort :: quick-sort (int a[], int l,
                               int h)
```

```
{
  int temp, key, low, high;
```

```
  low = l;
```

```
  high = h;
```

```
  key = a[(low + high) / 2];
```

```
  do
```

```
{
  while (key > a[low] → i.e. key > a[i]
        low stable)
  {
    low++;
```

```
  }
  while (key < a[high])
  {
    high--;
```

```
  }
  if (low <= high)
```

```
{
  temp = a[low];
  a[low + 1] = a[high];
  a[high - 1] = temp;
}
```

```
  }
  while (low <= high);
  if (l < high)
    quick-sort (a, l, high);
  if (low < h)
    quick-sort (a, low, h);
```

```
  }
  void quicksort :: output (int a[], int)
```

```
{
  for (i = 0; i <= n - 1; i++)
  {
    cout << endl << a[i];
```

```
  }
}
```

Q write a C program to implement Merge sort.

```
#include <stdio.h>
#include <conio.h>
void merge(int a[], int, int, int);
void merge-sort(int a[], int, int);
void main()
{
    int a[10], i, n;
    clrscr();
    printf("\n Enter the number of elements in the array:");
    scanf("%d", &n);
    printf("\n Enter the elements of the array:");
    for (i=0; i<n; i++)
    {
        scanf("%d", &a[i]);
    }
    merge-sort(a, 0, n-1);
    printf("\n The sorted array is:");
    for (i=0; i<n; i++)
        printf("%d\t", a[i]);
    getch();
}

void merge(int a[], int beg, int mid, int end)
{
    int i = beg, j = mid+1, index = beg; temp[10], k;
    while ((i <= mid) && (j <= end))
    {
        if (a[i] < a[j])
        {
            temp[index] = a[i];
            i++;
        }
    }
}
```

```

else
{
temp [index] = a[j];
j++;
}
index++;
}
if (i > mid)
{
while (j <= end)
{
temp[index] = a[j];
j++;
index++;
}
}
}

```

```

else
{
while (i <= mid)
{
temp[index] = a[i];
i++;
index++;
}
}
}

```

```

for (k = beg; k < index; k++)
a[k] = temp[k];
}

```

```

void merge-sort (int a[], int beg, int end)
{
int mid;
if (beg < end)
{
mid = (beg + end) / 2;
merge-sort (a, beg, mid);
merge-sort (a, mid + 1, end);
}
}

```


Heap sort

Algo Heapsort (ARR, N)

Step 1: [Build Heap]

Repeat for $I = 0$ to $N-1$

CALL Insert-Heap (ARR, N, ARR [I])

END OF LOOP

Step 2: (Repeatedly delete the root elem)

Repeat while $N > 0$

CALL Delete-Heap (ARR, N, VAL)

SET $N = N - 1$

END OF LOOP

Step 3: END

Various Operations of Heap

consider $\text{Heap-Size}(A)$ a variable containing the array index of the last element of the heap A . Elements in the array that are beyond this index are not part of heap. Thus the last element of heap is referenced by $A[\text{HEAP-SIZE}(A)]$

Traversal of Heap

parent:-

This operation returns index of parent

Algo :- parent operation

step I :- start

step II :- Return $i/2$

step 3 :- End

left:- This operation returns index of left child, let us assume that index i is given to us.
consider $\text{HEAP-SIZE}(A)$ a variable containing the array index of the last element of heap:-

Algo :- left operation

step ① start

step ② compare $\text{Heap-Size}(A)$ with $2i$ if $2i > \text{HEAP-SIZE}(A)$ then go to step ③ else go to

step ④

step ③ Return i is leaf node
step ④ Return $2i$
step ⑤ End

Algo Right operation

step ① start
step ② Compare Heap-size (A) with $2i+1$, if
 $2i+1 > \text{Heap-size (A)}$ then go to step ③
else go to step ④

step ③ Return i is leaf node
step ④ Return $2i+1$
step ⑤ End

Algo Heapify operation

step ① start
step ② Repeat the following steps for all
levels of Heap
step ③ compare the value of i with the
children of i (left & right). If the value
of i is the smallest of three then go to
step ⑤ else go to step ④
step ④ swap the position with the smaller
child.
step ⑤ End

Worst Case Comparisons

It takes 2 comparisons to move down 1 level i.e. 1 iteration of loop, comparing i to left child & then the smaller to move to i 's right child. The max no of times the procedure will take these comparisons is the height of the heap i.e. $\log_2 n$. Therefore the cost of heapify is approximately $2 \log_2 n$ where n 's the no of elements of heap rooted at i .

This worst case no of comparisons is $O(\log_2 n)$.

Build-Heap

Algo :- Build-Heap operation

step (1) start

step (2) starting from $i = n/4 * 2$ repeat step (3)

step (3) until it reaches the first element of the array.

step (4) call the operation heapify(i, A)

step (5) $i = i - 1$

step (6) $i = 1$

Worse Case Comparisons (Heap-Build) is better than Insert operators to build a heap.

The elements in the last level of heap are already single element heaps & therefore $\text{Build_Heap}(i, A)$ starts at the last element of the second last level & works its way back one element at a time until it reaches the first element of the array. $\text{Heapify}(i, A)$ uses two comparisons for each call.

Therefore at the second last level which contains $n/4$ elements, the cost will be $n/4 * 2$. For the third last level, which contains $n/8$ elements, the cost will be $n/8 * 4$, since each level is two comparisons. Summing the cost each level gives the cost of $\text{Build_Heap}(A)$ as: $\sim 2n = O(n)$. This is much better than using Insert(K) trees, $O(n \log_2 n)$.

Analysis of sorting Method

Sorting	Best Case	Average Case	Worst Case
Bubble	$O(n^2)$	$O(n^2)$	$O(n^2)$
Insertion	$O(n)$	$O(n^2)$	$O(n^2)$
Shell	$O(n)$	$O(n^2)$	$O(n^2)$
Heap	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Merge	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Quick	$O(\log n)$	$O(\log n)$	$O(n^2)$

Internal sorting :- This sorting method is applied on the data stored on the main memory. Since the main memory is not feasible to apply the sorting technique always

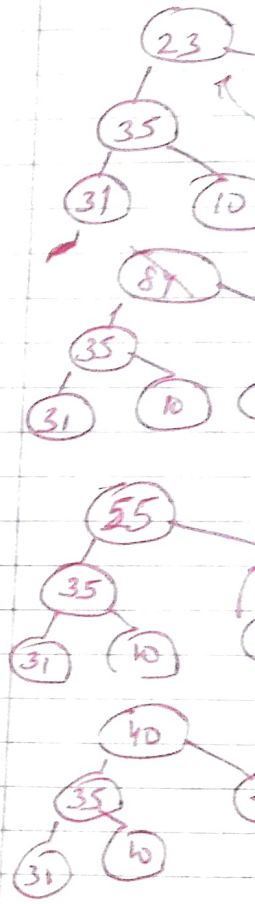
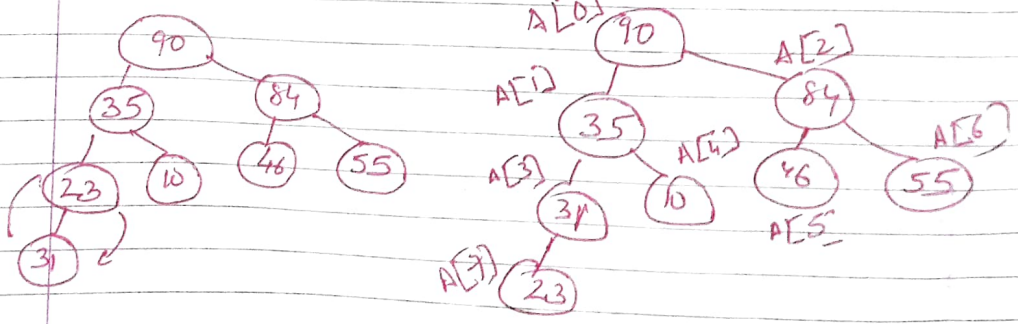
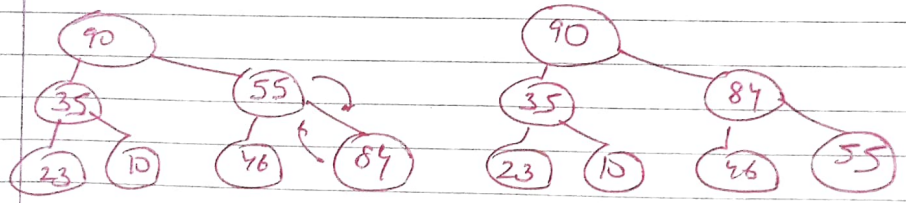
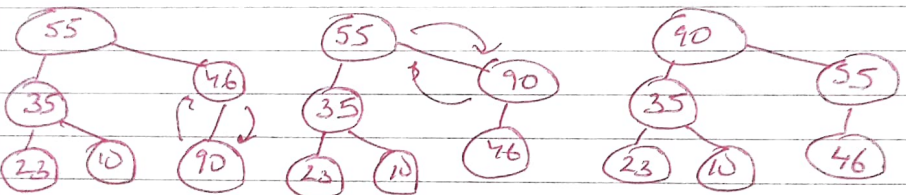
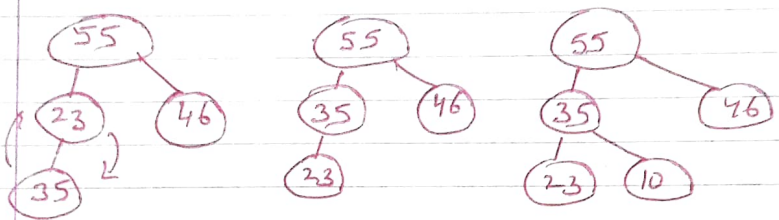
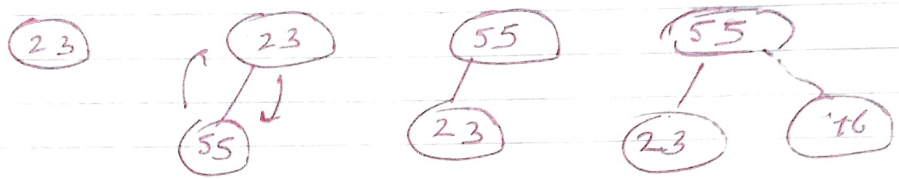
External sorting

External sorting is a complex sorting technique when a large data has to be sorted & if we need to place part of data on the main memory & remaining on the secondary memory then external sorting can be done.

The data stored on secondary memory is part by part loaded into the main memory sorting can be done over there. The sorted data can be then stored in the intermediate files. Thus many intermediate files can be generated. And finally all the sorted intermediate files are merged in a single file. Thus the huge amount of data can be sorted using this technique.

23 55 46 35 10 90 8 31

deletes



Deletes 90 & insert 23 at top



Q Write a C program to implement heap sort.

```
#include <stdio.h>
#include <conio.h>
#define MAX 10
void RestoreHeapUp(int *, int);
void RestoreHeapDown(int *, int, int);
void main()
{
    int Heap[MAX], n, i, j;
    clrscr();
    printf("\n Enter the number of elements:");
    scanf("%d", &n);
    printf("\n Enter the elements");
    for (i = 1; i <= n; i++)
    {
        scanf("%d", &Heap[i]);
        RestoreHeapUp(Heap, i);
    }
    // Delete the root element & heapify the heap
    j = n;
    for (i = 1; i <= j; i++)
    {
        int temp;
        temp = Heap[1];
        Heap[1] = Heap[n];
        Heap[n] = temp;
        n = n - 1; // The elem Heap[n] to be deleted suppress
        RestoreHeapDown(Heap, 1, n); // heapify
    }
    n = j;
    printf("\n The sorted elements are:");
}
```